MPI types, Scatter and Scatterv

Version 2

MPI types, Scatter and Scatterv

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

Logical and physical layout of a C/C++ array in memory.

A = malloc(6*6*sizeof(int));

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35

MPI_Scatter

```
int MPI_Scatter(
    const void *sendbuf,   // data to send
    int sendcount,         // sent to each process
    MPI_Datatype sendtype,// type of data sent
    void *recvbuf,         // where received
    int recvcount,         // how much to receive
    MPI_Datatype recvtype,// type of data received
    int root,              // sending process
    MPI_Comm comm)         // communicator
```

`sendbuf`, `sendcount`, `sendtype` valid only at the sending process

Equal number elements to all processors

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
int MPI_Scatter(A, 9, MPI_Int, B, 9,
                MPI_Int,0,
                MPI_COMM_WORLD)
```

P_0	0	1	2	3	4	5	6	7	8
P_1	9	10	11	12	13	14	15	16	17
P_2	18	19	20	21	22	23	24	25	26
P_3	27	28	29	30	31	32	33	34	35

MPI_Scatterv

```
int MPI_Scatter(
    const void *sendbuf,    // data to send
    const int *sendcounts,  // sent to each process
    const int* displ        // where in sendbuf
                            // sent data is
    MPI_Datatype sendtype,  // type of data sent
    void *recvbuf,          // where received
    int recvcount,          // how much to receive
    MPI_Datatype recvtype,  // type of data received
    int root,               // sending process
    MPI_Comm comm)          // communicator
```

`sendbuf`, `sendcount`, `sendtype` valid only at the sending process

Specify the number elements sent to each processor

```
int[] counts = {10, 9, 8, 9};
int[] displ = {0, 10, 19, 27};
int MPI_Scatterv(A, counts, displs, MPI_Int,rb, counts, MPI_Int 0,
                 MPI_COMM_WORLD)
```

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

rb

P_0	0	1	2	3	4	5	6	7	8	9
P_1	10	11	12	13	14	15	16	17	18	
P_2	19	20	21	22	23	24	25	26		
P_3	27	28	29	30	31	32	33	34	35	

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35

MPI_Type_vector

```
int MPI_Type_vector(
      int count,              // number of blocks
      int blocklength,        // #elts in a block
      int stride,             // #elts between block starts
      MPI_Datatype oldtype,   // type of block elements
      MPI_Datatype *newtype   // handle for new type
)
```

Allows a type to be created that puts together blocks of elements in a vector into another vector.

Note that a 2-D array in contiguous memory can be treated as a 1-D vector.

MPI_Type_vector: defining the type

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype col, coltype;
MPI_Type_vector(6, 1, 6, MPI_INT,
                &col);
MPI_Type_commit(&col);
MPI_Send(A, 1, col, P-1,
         MPI_ANY_TAG,
         MPI_Comm_World);
```

There are 6 blocks, and each is made of 1 int, and the new block starts 6 positions in the linearized array from the start of the previous block.

0	6	12	18	24	30

MPI_Type_vector: defining the type

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

blocks

elts in a block

elts between block starts

```
MPI_Datatype col, coltype;
MPI_Type_vector(6, 1, 6, MPI_INT,
                &col);
MPI_Type_commit(&col);
MPI_Send(A, 1, col, P-1,
         MPI_ANY_TAG,
         MPI_Comm_World);
```

There are 6 blocks, and each is made of 1 int, and the new block starts 6 positions in the linearized array from the start of the previous block.

0	6	12	18	24	30

What if we want to scatter columns (C array layout)

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

P_0	0	6	12	18	24	30
P_1	1	7	13	19	25	31
P_2	2	8	14	20	26	32
P_3	3	9	15	21	27	33
P_4	4	10	16	22	28	34
P_5	5	11	17	23	29	35

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35

What if we want to scatter columns?

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

blocks *(note change from handouts)*

elts in a block

elts between block starts

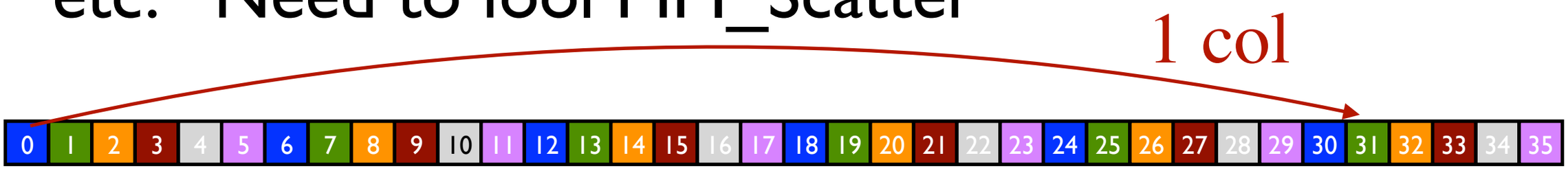
```
MPI_Datatype col, coltype;
MPI_Type_vector(6, 1, 6, MPI_INT,
                &col);
MPI_Type_commit(&col);
int MPI_Scatter(A, 6, col, AC, 6,
                MPI_Int, 0,
                MPI_Comm_World);
```

The code above won't work.

Why?

Where does the first col end?

We want the first column to end at 0, the second at 1, etc. Need to fool MPI_Scatter



MPI_Type_create_resized to the rescue

```
int MPI_Type_create_resized(
    MPI_Datatype oldtype, // type being resized
    MPI_Aint lb, // new lower bound
    MPI_Aint extent, // new extent (“length”)
    MPI_Datatype *newtype) // resized type name
)
```

Allows a new size (or *extent*) to be assigned to an existing type.

Allows MPI to determine how far from an object O1 the next adjacent object O2 is. As we will see this is often necessitated because we treat a logically 2-D array as a 1-D vector.

Using MPI_Type_vector

A


```
MPI_Datatype col, coltype;

MPI_Type_vector(6, 1, 6, MPI_INT,&col);

MPI_Type_commit(&col);

MPI_Type_create_resized(col, 0,
1*sizeof(int), &coltype);

MPI_Type_commit(&coltype);

MPI_Scatter(A, 1, coltype, rb, 6,
MPI_Int, 0, MPI_COMM_WORLD);
```

Using MPI_Type_vector

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

New length (extent)

New type

Old type

New lower bound

```
MPI_Datatype col, coltype

MPI_Type_vector(6, 1, 6, MPI_INT,&col);

MPI_Type_commit(&col);

MPI_Type_create_resized(col, 0,
1*sizeof(int), &coltype);

MPI_Type_commit(&coltype);

MPI_Scatter(A, 1, coltype, rb, 6,
MPI_Int, 0, MPI_COMM_WORLD);
```

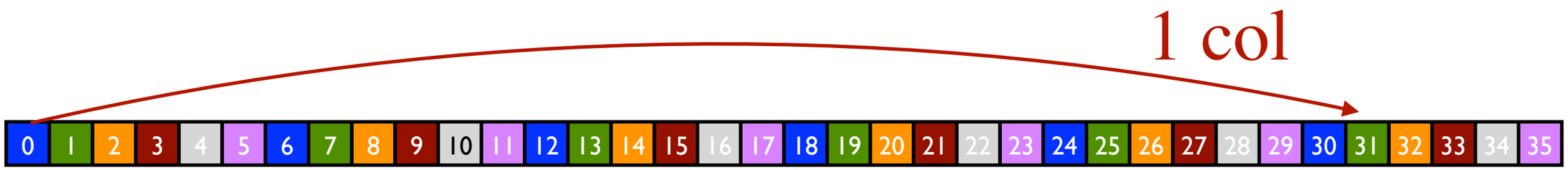
MPI_Type_vector: defining the type

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype col, coltype;
MPI_Type_vector(6, 1, 6, MPI_INT,
                        &col);
MPI_Type_commit(&col);
MPI_Type_create_resized(col, 0,
    1*sizeof(int), &coltype);
MPI_Type_commit(&coltype);
MPI_Scatter(A, 1, coltype, rb,
    6, MPI_Int, 0, MPI_COMM_WORLD);
```

Again, there are 6 blocks, and each is made of 1 int, and the new block starts 6 positions in the linearized array from the start of the previous block.



Using MPI_type_create_resized

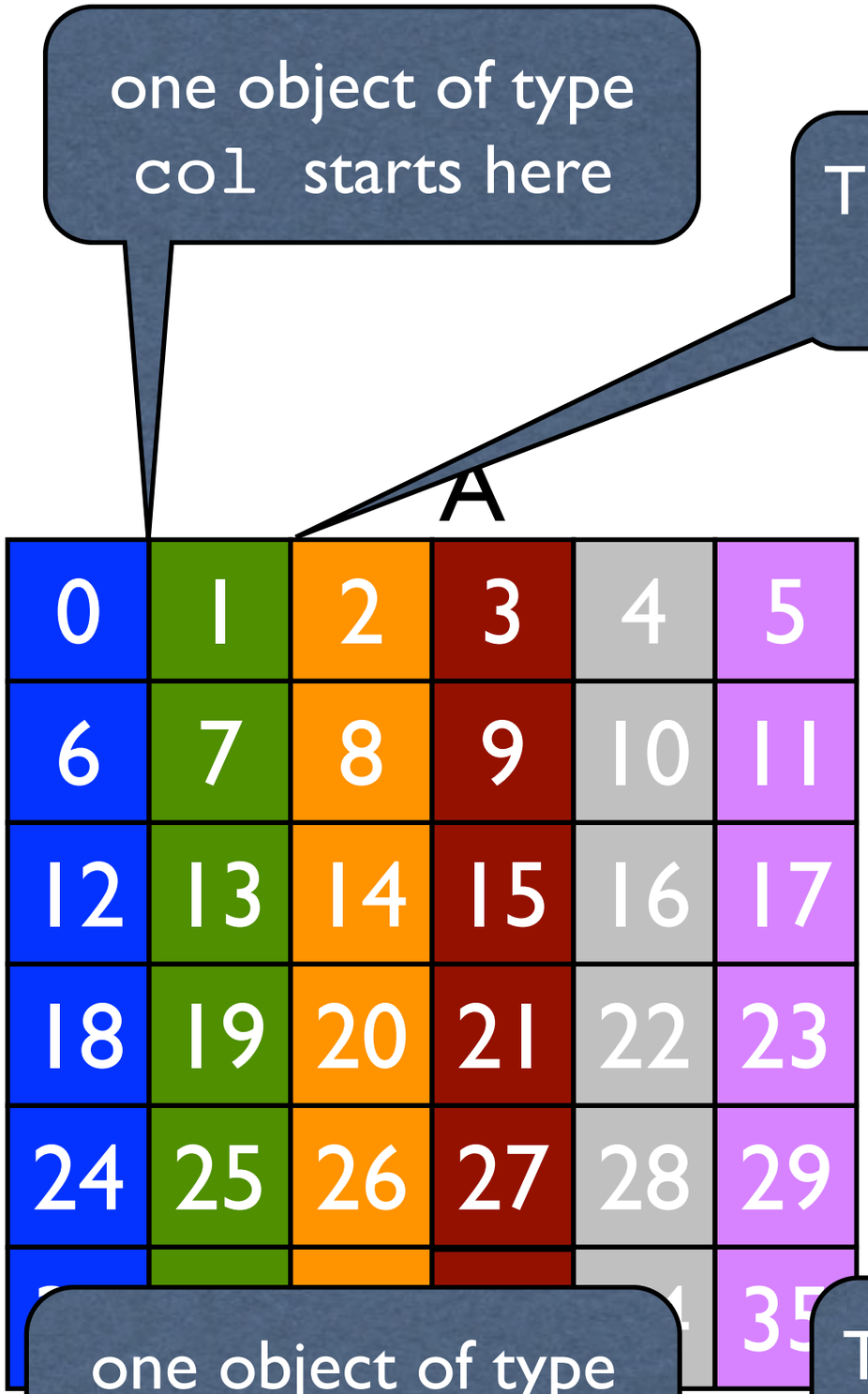
A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype col, coltype;
MPI_Type_vector(6, 1, 6, MPI_INT,
                &col);
MPI_Type_commit(&col);
MPI_Type_create_resized(col, 0,
    1*sizeof(int), &coltype);
MPI_Type_commit(&coltype);
MPI_Scatter(A, 1, coltype, rb,
    6, MPI_Int, 0, MPI_COMM_WORLD);
```

resize creates a new type from a previous type and changes the *size*. This allows easier computation of the offset from one element of a type to the next element of a type in the original data structure.

1 word, i.e. 1 coltype




```
MPI_Datatype col, coltype;
MPI_Type_vector(6, 1, 6, MPI_INT,
                    &col);
MPI_Type_commit(&col);
MPI_Type_create_resized(col, 0,
    1*sizeof(int), &coltype);
MPI_Type_commit(&coltype);
MPI_Scatter(A, 1, coltype, rb,
    6, MPI_Int, 0, MPI_COMM_WORLD);
```

resize creates a new type from a previous type and changes the *size*. This allows

of the offset from one

the next element of

data structure.

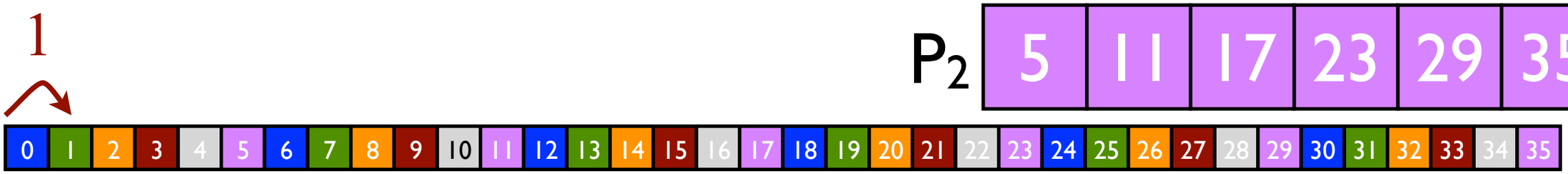
The result of the communication

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype col, coltype;
MPI_Type_vector(6, 1, 6, MPI_INT,
                &col);
MPI_Type_commit(&col);
MPI_Type_create_resized(col, 0,
    1*sizeof(int), &coltype);
MPI_Type_commit(&coltype);
MPI_Scatter(A, 1, coltype, rb,
    6, MPI_Int, 0, MPI_COMM_WORLD);
```

P_0	0	6	12	18	24	30
P_1	1	7	13	19	25	31
...						
P_2	5	11	17	23	29	35



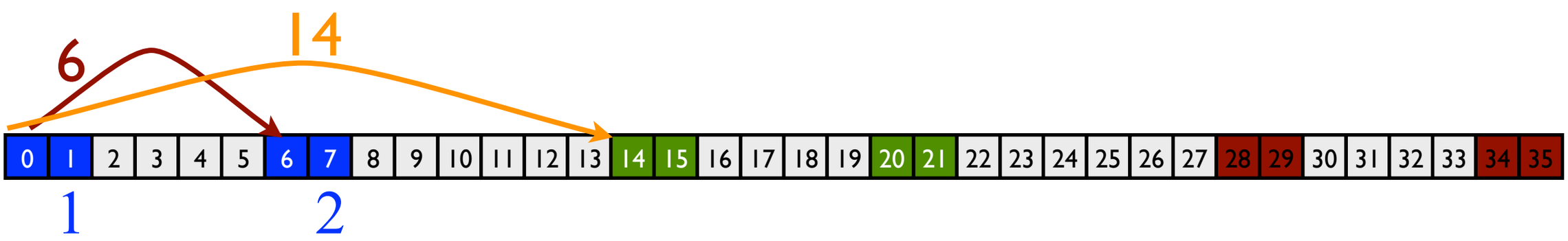
Scattering diagonal blocks

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype block, blocktype;
MPI_Type_vector(2, 2, 6, MPI_INT,
                    &block);
MPI_Type_commit(&block);
MPI_Type_create_resized(block, 0,
    14*sizeof(int), &blocktype);
MPI_Type_commit(&blocktype);
int MPI_Scatter(A, 1, blocktype, B, 4,
                    MPI_Int,0,
                    MPI_COMM_WORLD)
```

note that 2*numrows + width of block =14



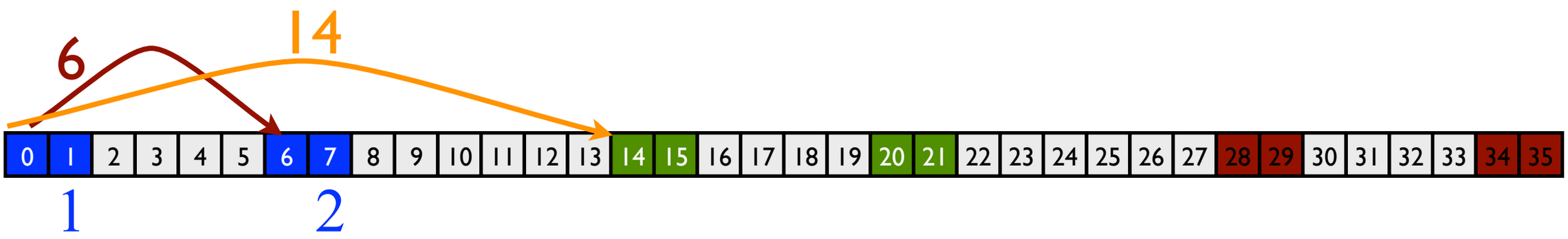
Scattering diagonal blocks

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype block, blocktype;
MPI_Type_vector(2, 2, 6, MPI_INT,
                    &block);
MPI_Type_commit(&block);
MPI_Type_create_resized(block, 0,
    14*sizeof(int), &blocktype);
MPI_Type_commit(&blocktype);
int MPI_Scatter(A, 1, blocktype, B, 4,
                    MPI_Int,0,
                    MPI_COMM_WORLD)
```

note that 2*numrows + width of block =14 as shown by the red arrows.



Scattering the blocks

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype block, blocktype;
MPI_Type_vector(2, 2, 6, MPI_INT,
                    &block);
MPI_Type_commit(&block);
MPI_Type_create_resized(block, 0,
    14*sizeof(int), &blocktype);
MPI_Type_commit(&blocktype);
int MPI_Scatter(A, 1, blocktype, B, 4,
                    MPI_Int,0,
                    MPI_COMM_WORLD)
```

B

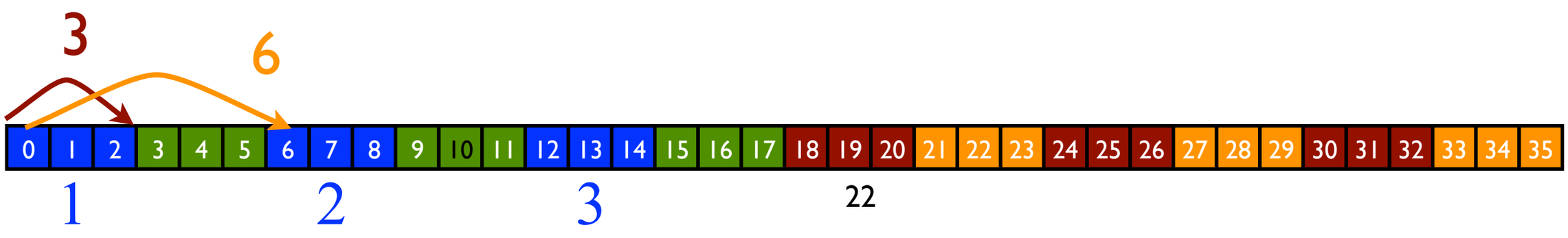
P_0	0	1	6	7
P_1	14	15	20	21
P_2	28	29	34	35

The Type_vector statement describing this

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype block, blocktype;
MPI_Type_vector(3, 3, 6, MPI_INT,
                        &block);
MPI_Type_commit(&block);
MPI_Type_create_resized(block, 0,
    3*sizeof(int), &blocktype);
MPI_Type_commit(&blocktype);
```



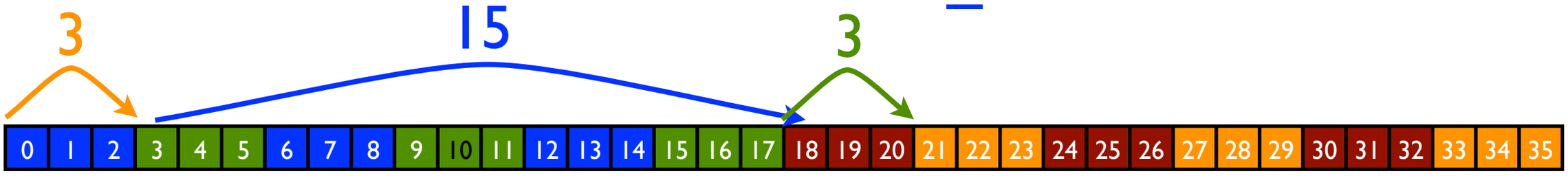
The create_resize statement for this

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

```
MPI_Datatype block, blocktype;
MPI_Type_vector(3, 3, 6, MPI_INT,
                    &block);
MPI_Type_commit(&block);
MPI_Type_create_resized(block, 0,
     3*sizeof(int), &blocktype);
MPI_Type_commit(&blocktype);
```

Distance between start of blocks varies, but are multiples of 3. Use MPI_Scatterv

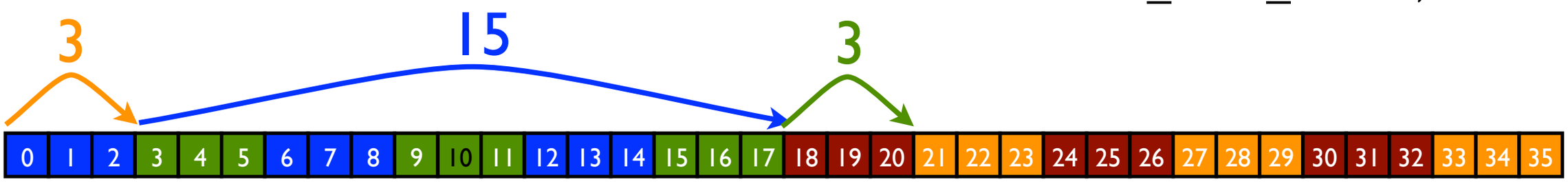


Sending the data

A

0	1	2	3	4	5
6	7	8	9	10	11
12	13	14	15	16	17
18	19	20	21	22	23
24	25	26	27	28	29
30	31	32	33	34	35

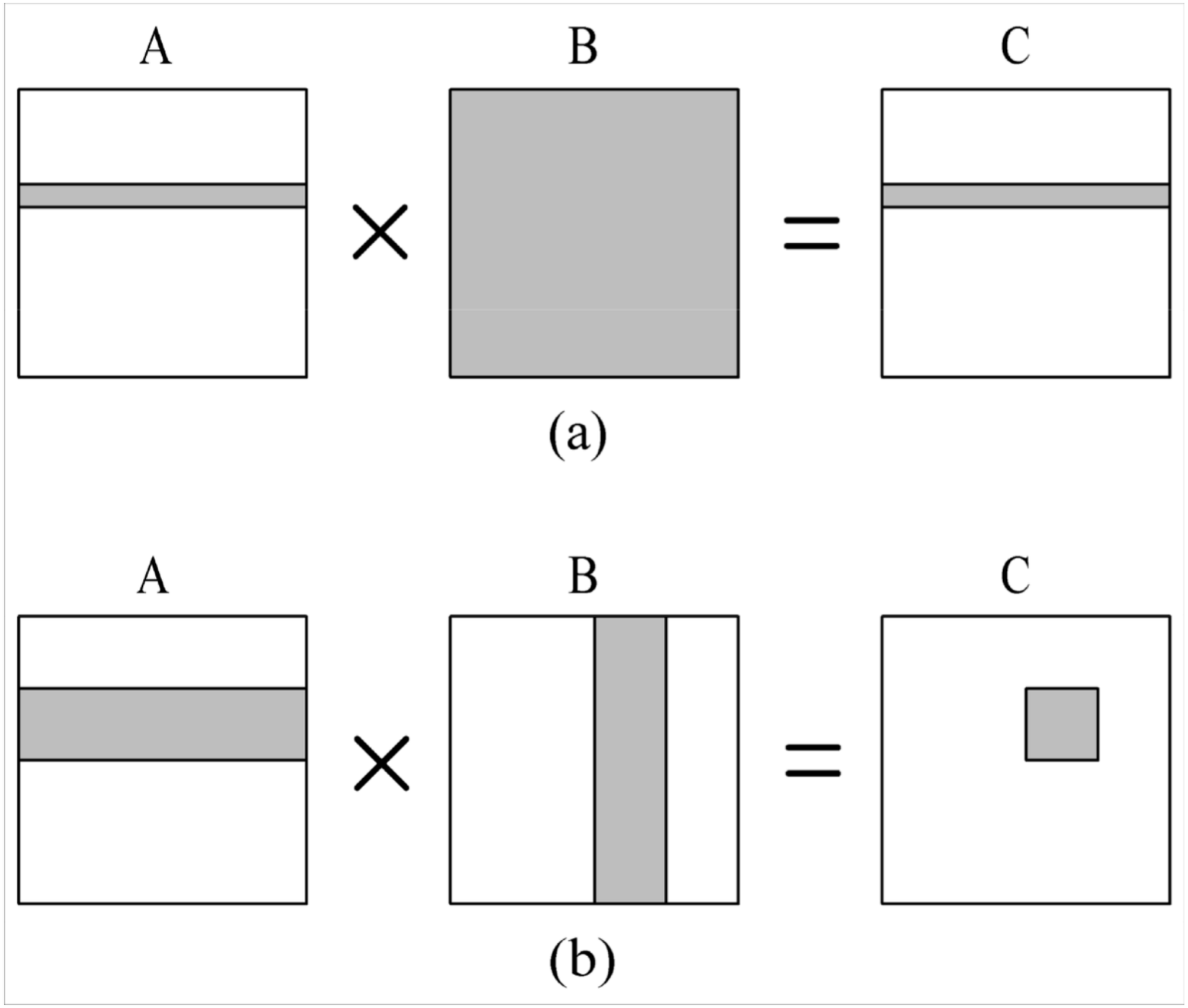
```
MPI_Datatype block, blocktype;
int disp = {0, 1, 6, 7)
int scount = {1, 1, 1, 1}
int rcount = {9, 9, 9, 9}
MPI_Type_vector(3, 3, 6, MPI_INT,
                        &block);
MPI_Type_commit(&block);
MPI_Type_create_resized(block, 0,
     3*sizeof(int), &blocktype);
MPI_Type_commit(&blocktype);
int MPI_Scatterv(A, scount, displ,
                    blocktype, rb, rcount,
                    MPI_Int, 0,
                    MPI_COMM_WORLD)
```



Matrix Multiply
Cannon's Algorithm

- Useful for the small project (ignore this!)
- Algorithm 1 in what follows is the layout we discussed earlier

Elements of A and B Needed to Compute a Process's Portion of C



Algorithm 1

Cannon's Algorithm

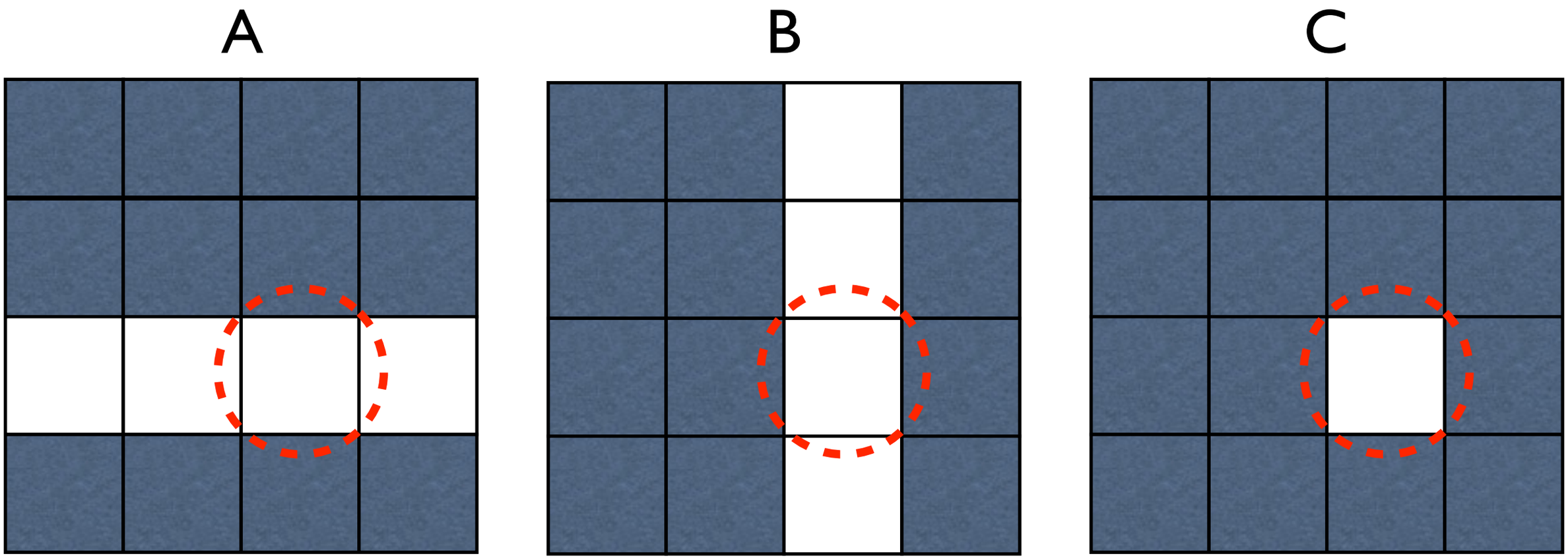
Parallel Algorithm 2 (Cannon's Algorithm)

- Associate a primitive task with each matrix element
- Agglomerate tasks responsible for a square (or nearly square) block of C (the result matrix)
- Computation-to-communication ratio rises to $n / \sqrt{p}$ (same total computation, more computation per communication) $2n / p < n / \sqrt{p}$ when p > 4

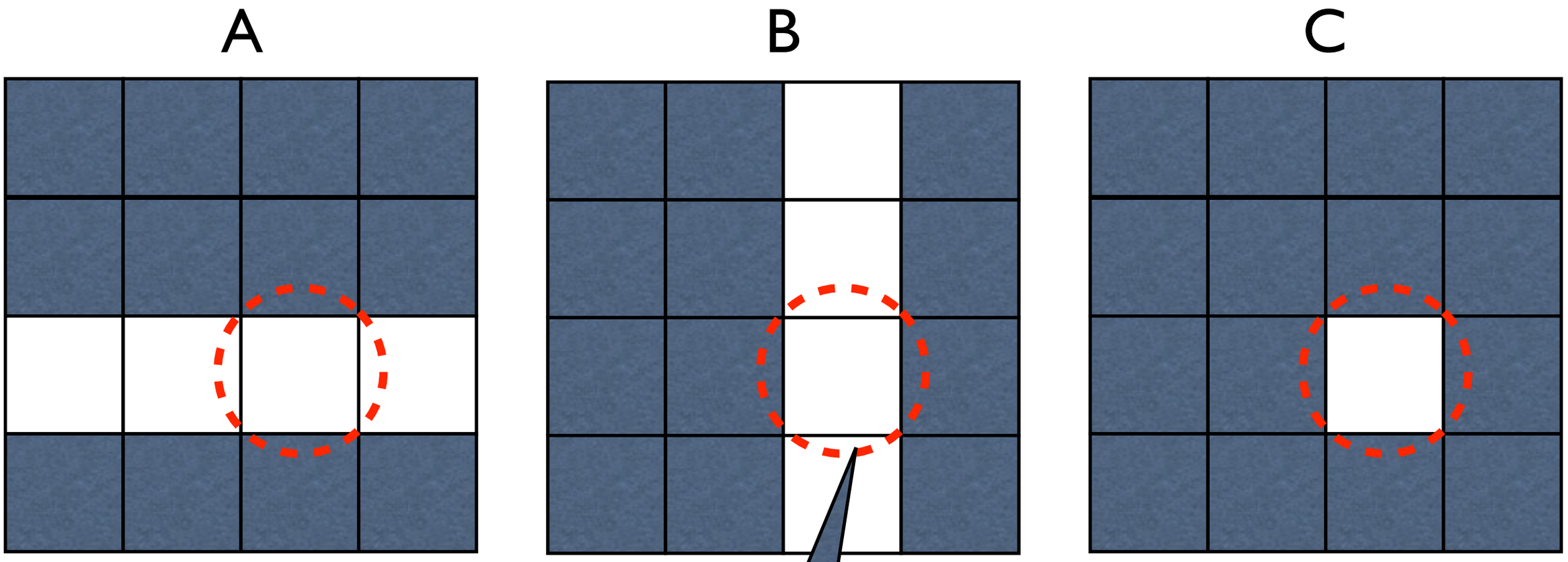
A simplifying assumption

- Assume that
  - *A*, *B* and (consequently) *C* are $n \times n$ square matrices
  - $\sqrt{p}$ is an integer, and
  - $n = k \cdot \sqrt{p}$, k an integer (i.e. n is a multiple of $\sqrt{p}$

Blocks need to compute part of a C element



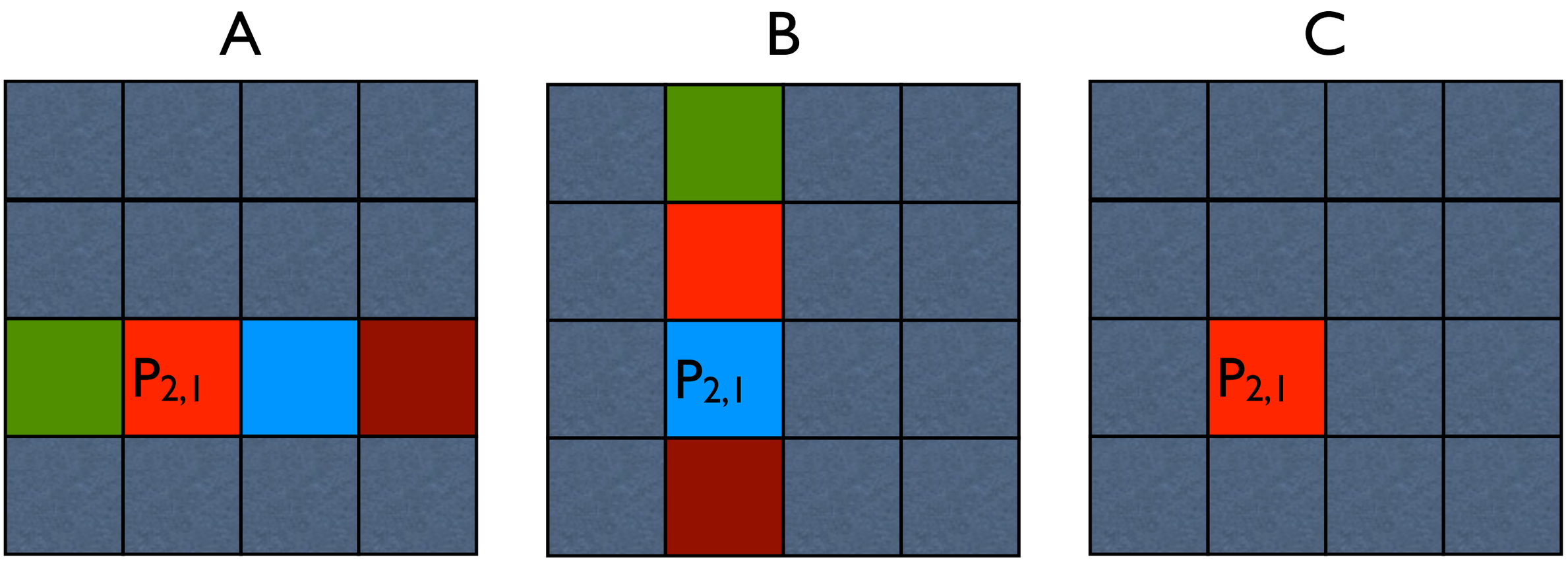
Blocks need to compute a C element



Blocks needed to compute a C element

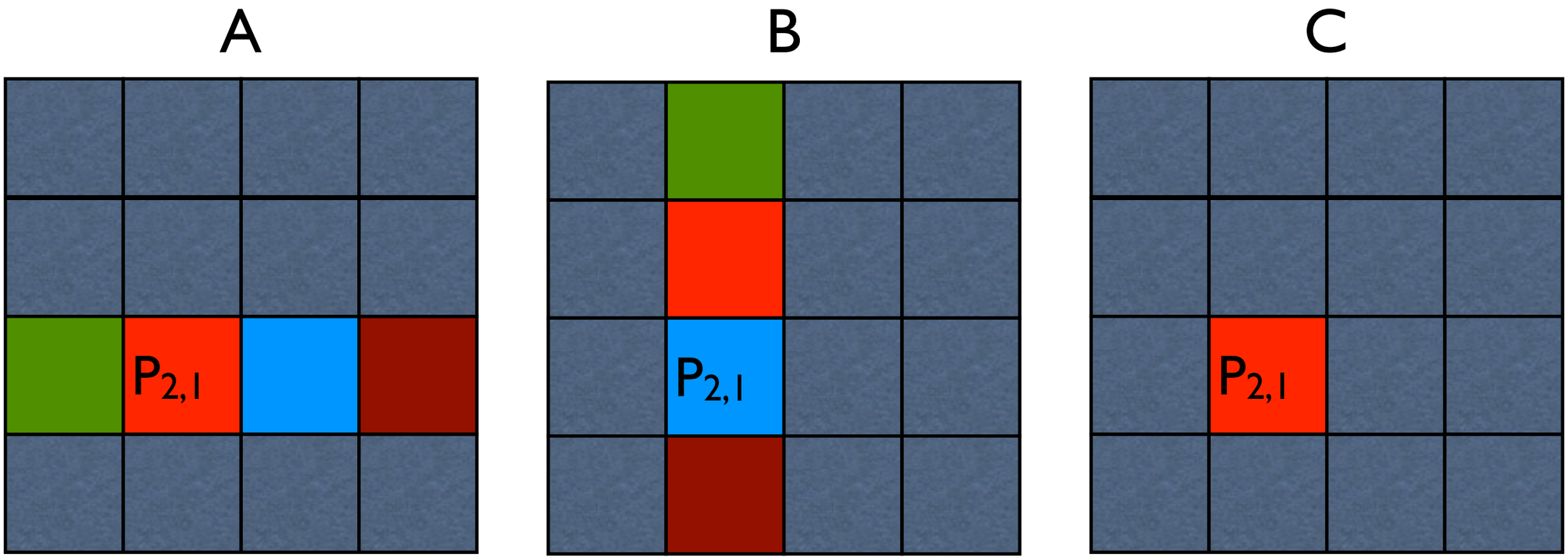
Processor $P_{2,1}$ needs, at some point, to simultaneously hold the green A and B blocks, the red A and B blocks, the blue A and B blocks, and the cayenne A and B blocks.

With the current data layout it cannot do useful work because it does not contain matching A and B blocks (it has a red A and blue B block)



Blocks needed to compute a C element

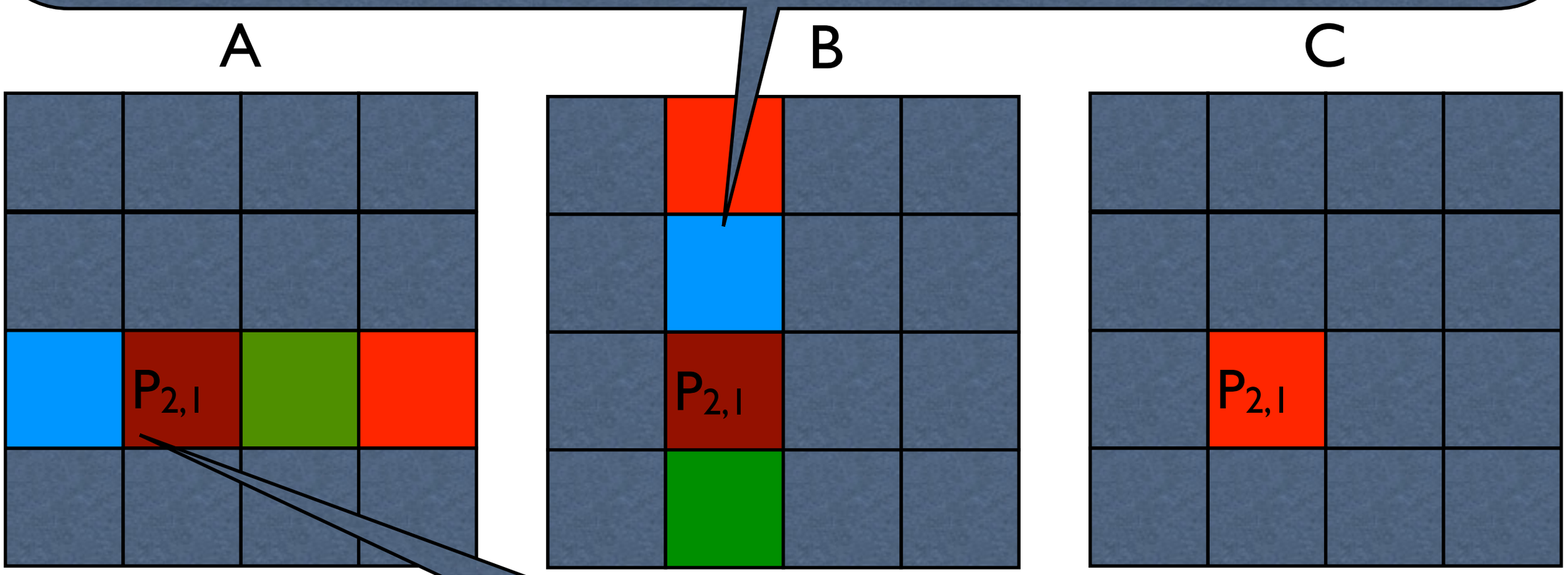
We need to rearrange the data so that every block has useful work to do



The initial data configuration does not provide for this

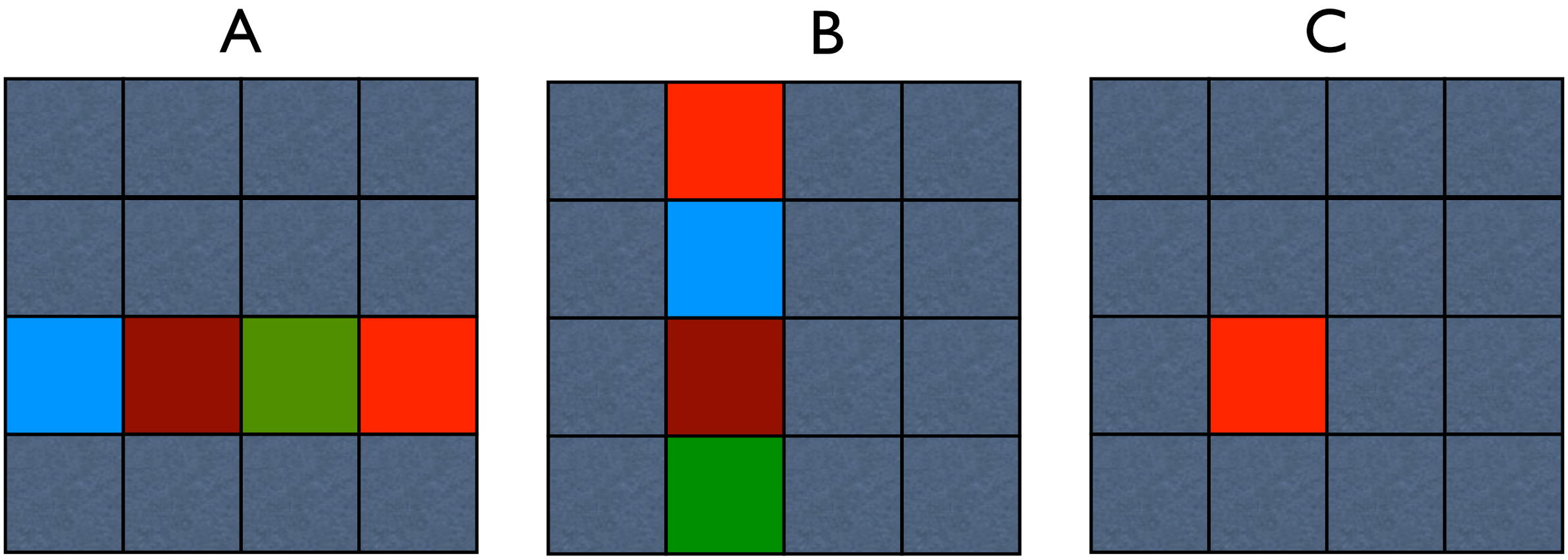
Change the initial data setup

Move every element $B_{i,j}$ up j rows



Move every element $A_{i,j}$ over i columns

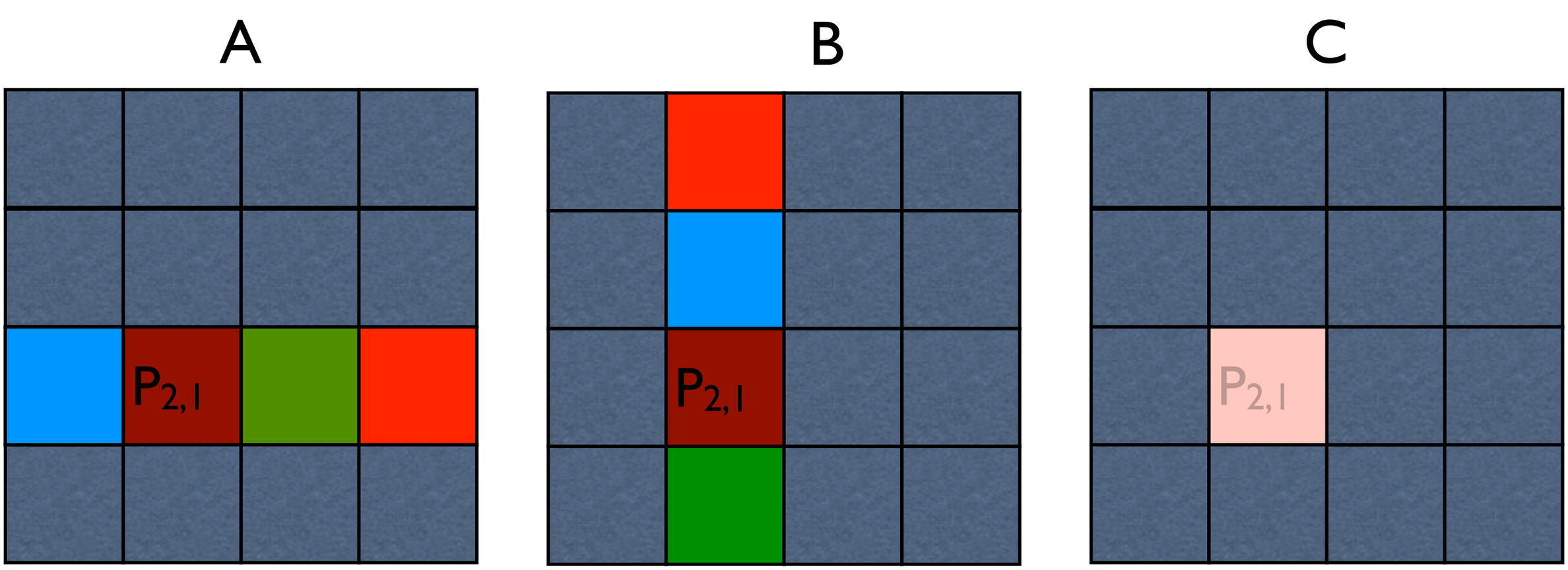
Every processor now has useful work to do



Note -- this only shows the full data layout for one processor

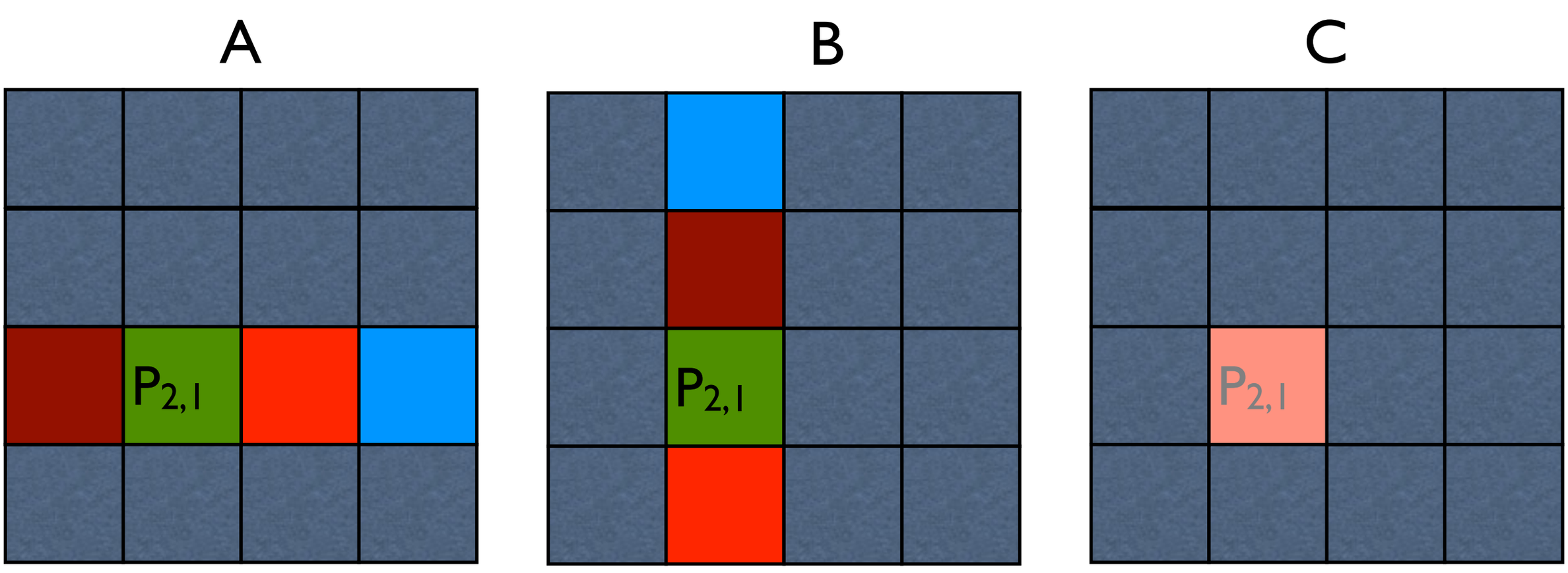
At each step in the multiplication, shift B elements up within their column, and A elements left within their row

First partial sum



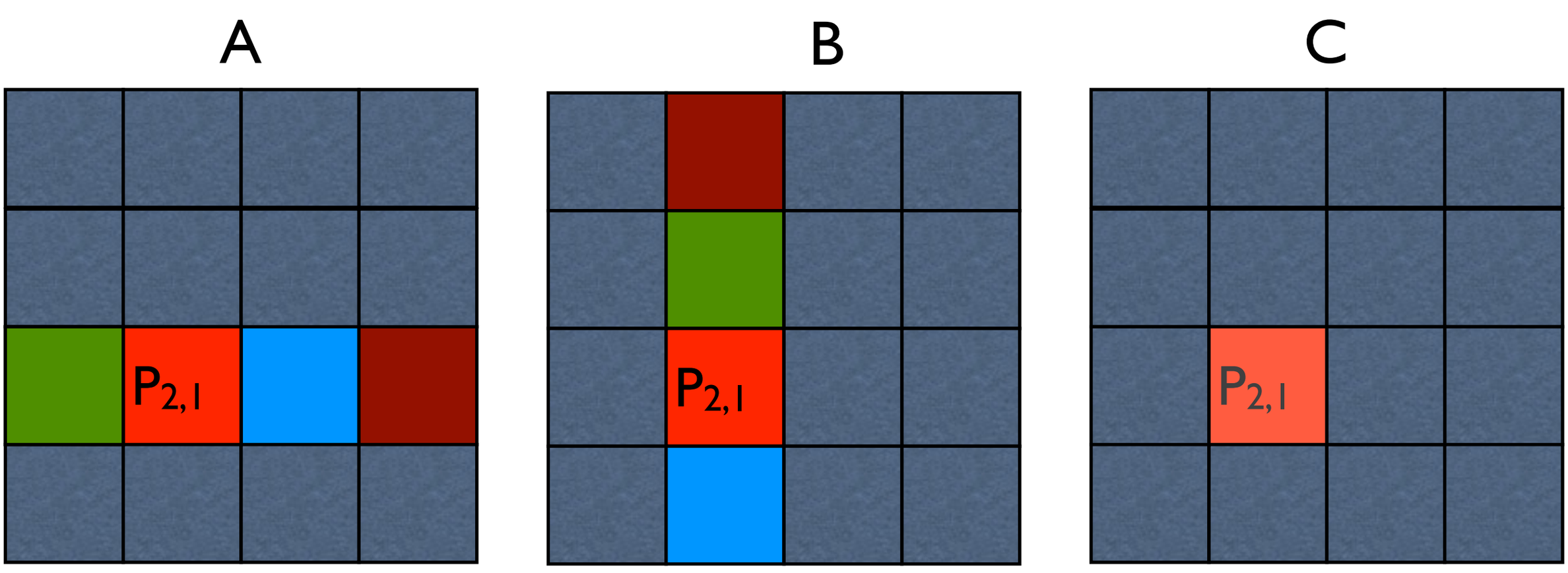
At each step in the multiplication, shift B elements up within their column, and A elements left within their row

Second partial sum



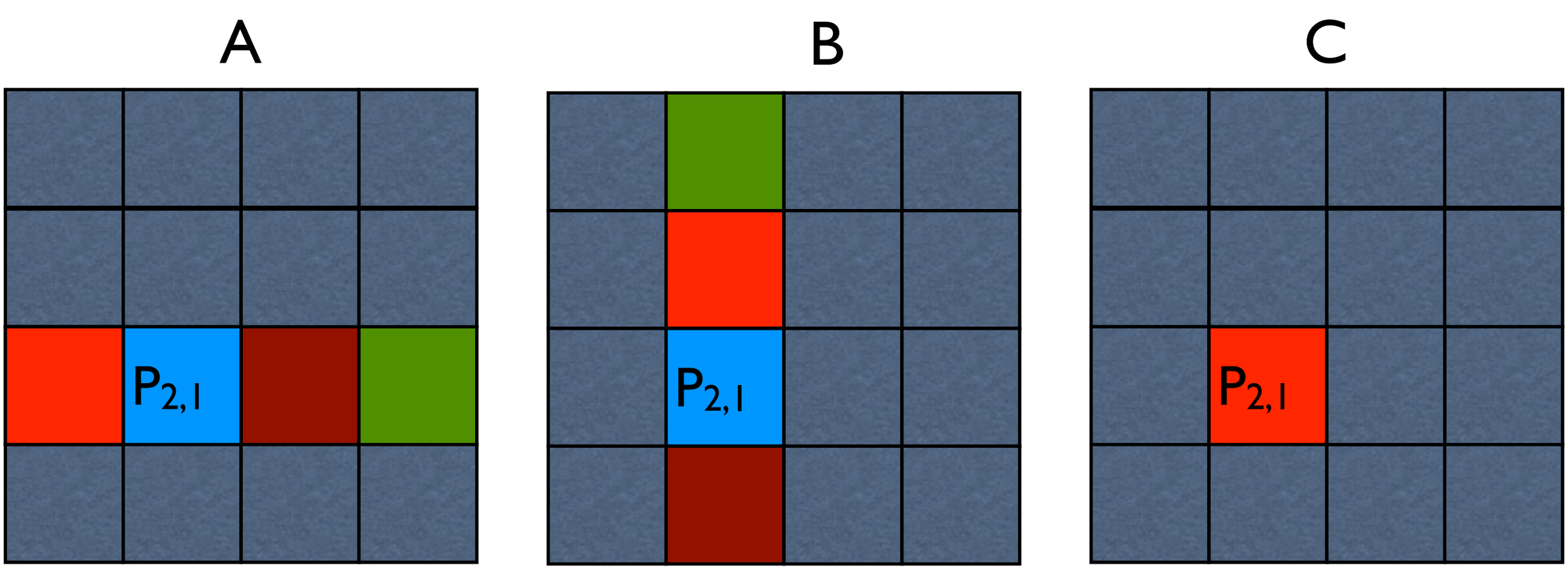
At each step in the multiplication, shift B elements up within their column, and A elements left within their row

Third partial sum

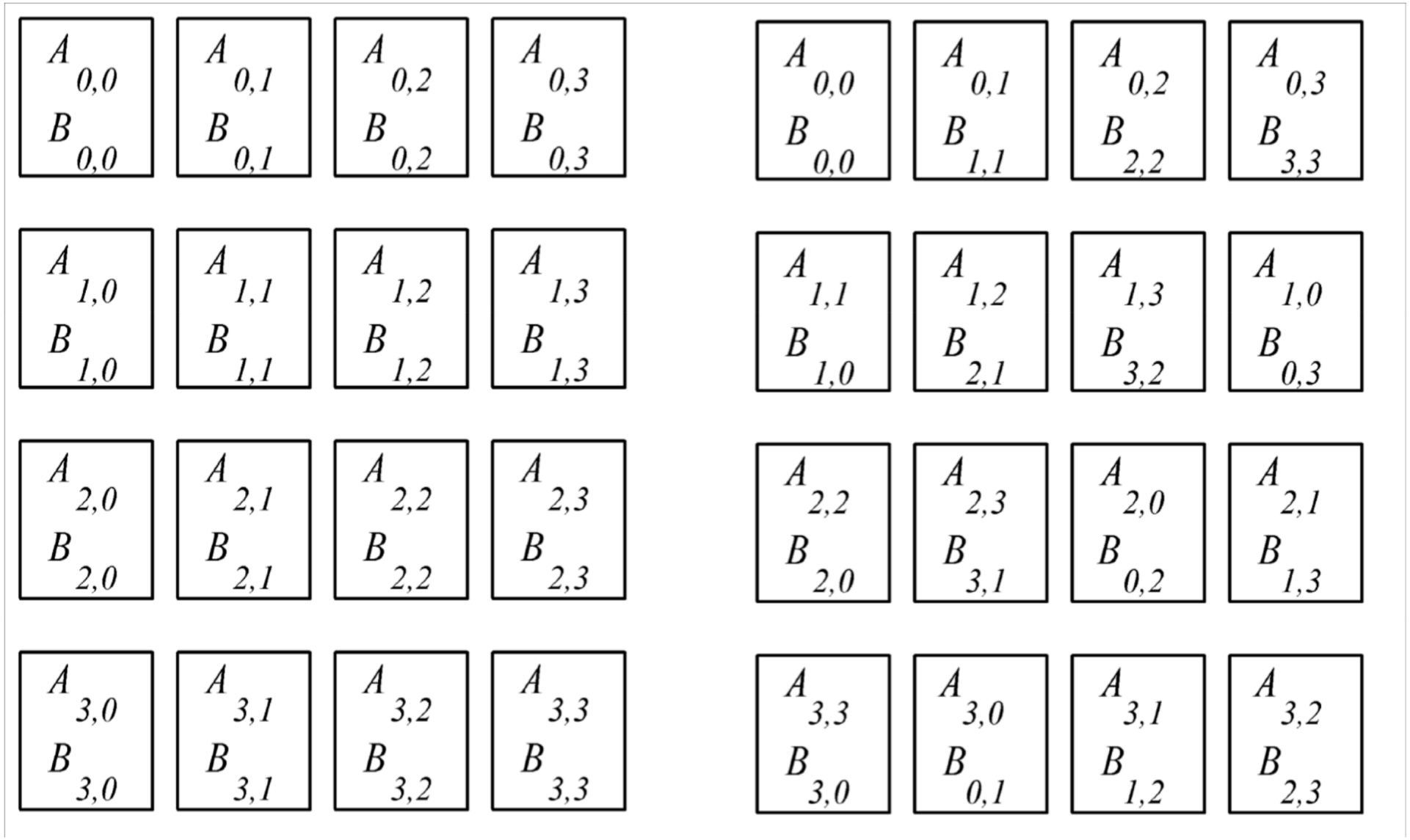


At each step in the multiplication, shift B elements up within their column, and A elements left within their row

Fourth partial sum



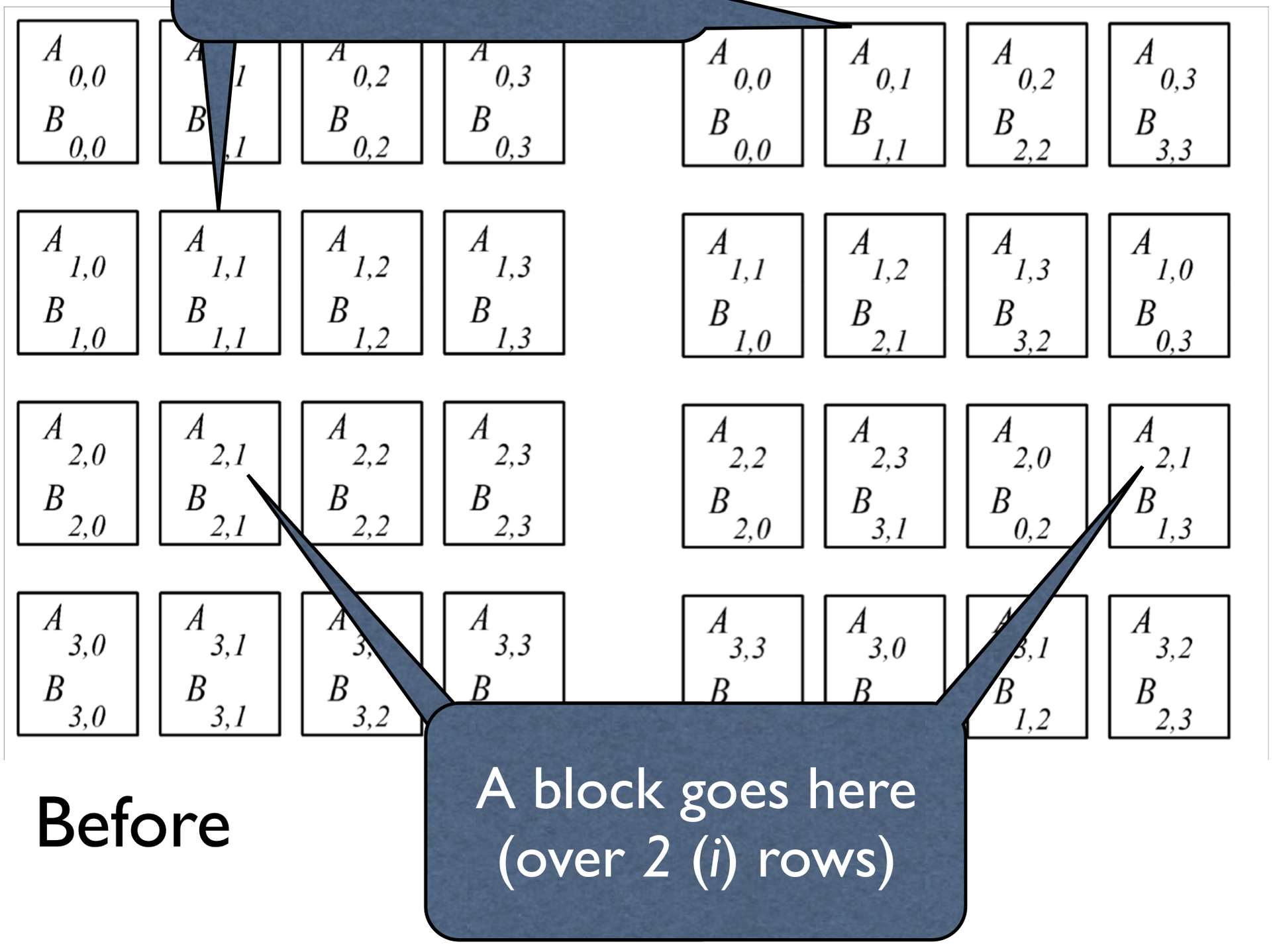
Another way to view this



Before

After

Anot o view this



Before

Yet another way to view this

Each triangle represents a matrix block on a processor

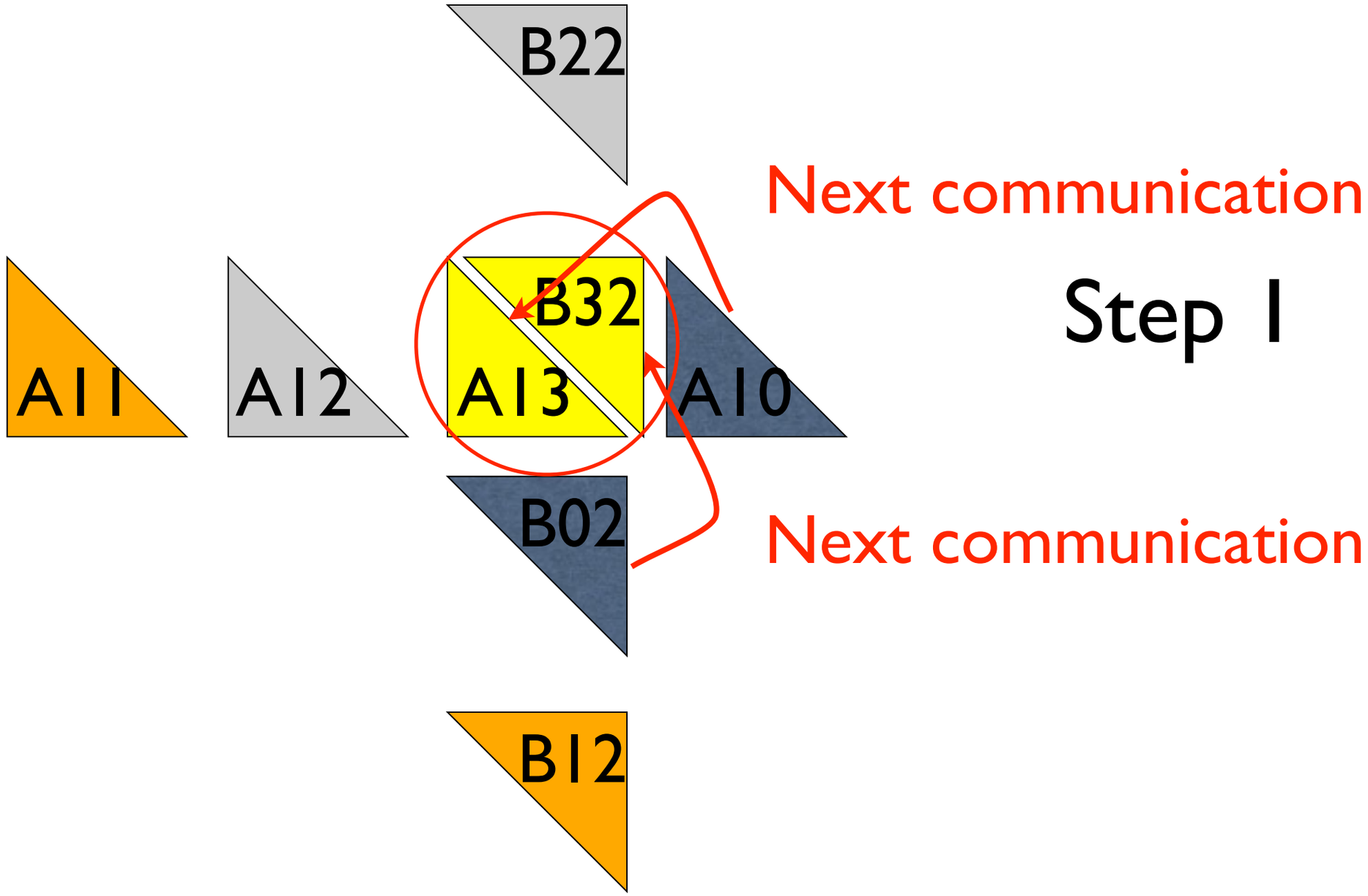
Only same-color triangles should be multiplied

Rearrange Blocks

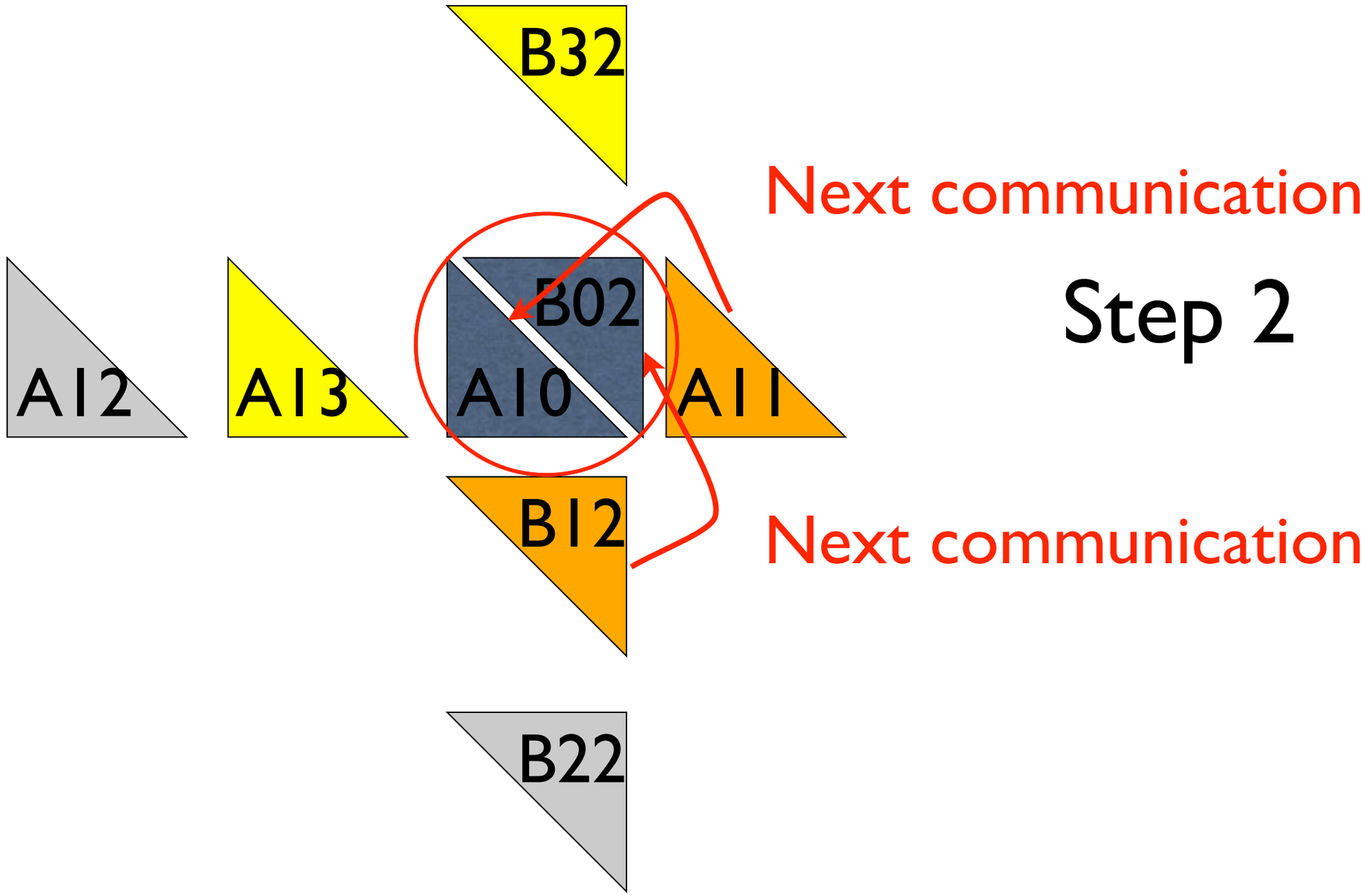
Block A*i,j* shifts left *i* positions

Block B*i,j* shifts up *j* positions

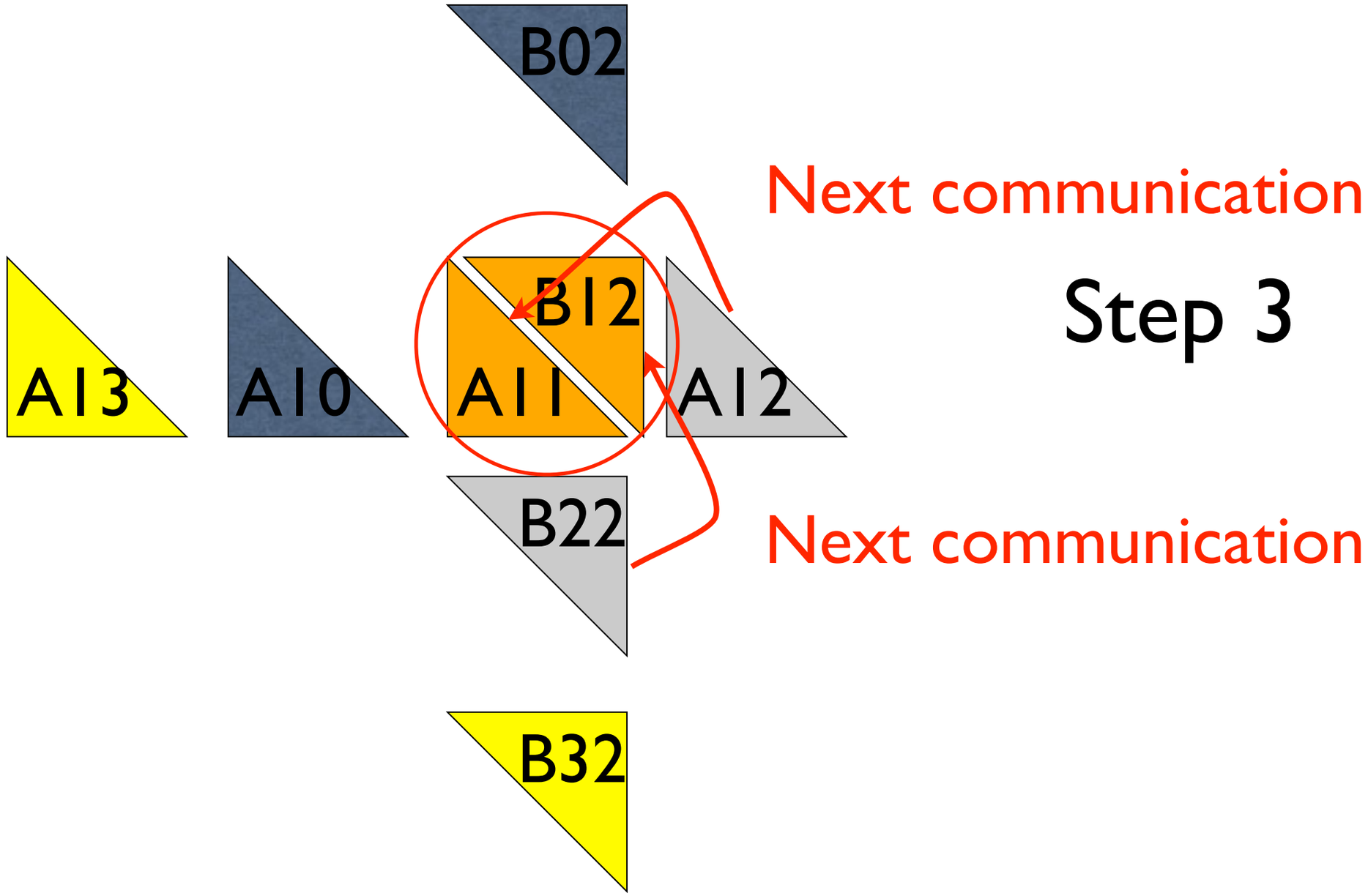
Consider Process $P_{1,2}$



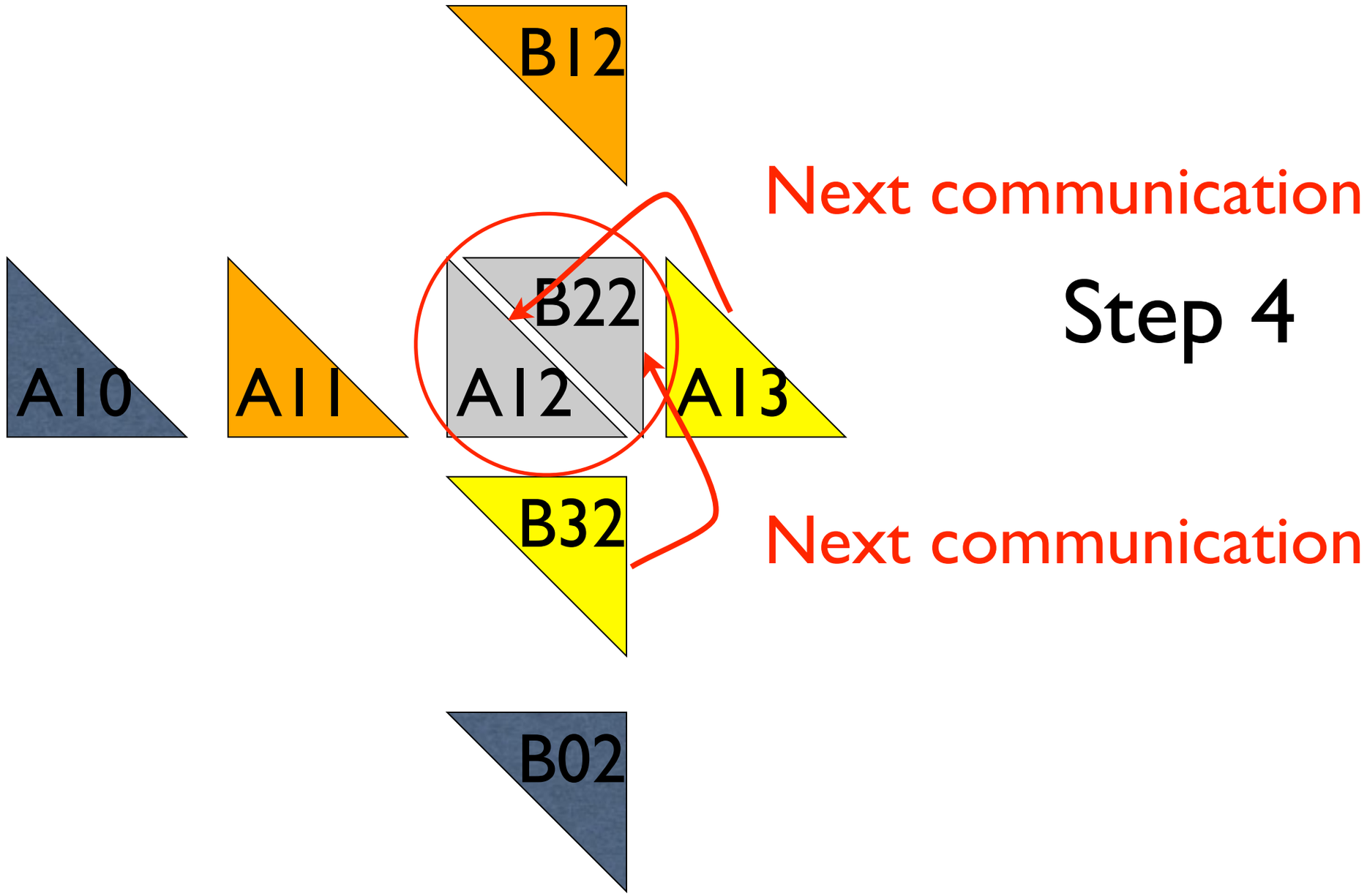
Consider Process $P_{1,2}$



Consider Process $P_{1,2}$



Consider Process $P_{1,2}$



Complexity Analysis

- Algorithm has $\sqrt{p}$ iterations
  - During each iteration process multiplies two $(n / \sqrt{p}) \times (n / \sqrt{p})$ matrices: $\Theta(n / \sqrt{p})^3$ or $\Theta(n^3 / p^{3/2})$
- Overall computational complexity: $\sqrt{p}\, n^3/p^{3/2}$ or $\Theta(n^3 / p)$
  - During each $\sqrt{p}$ iterations a process sends and receives two blocks of size $(n / \sqrt{p}) \times (n / \sqrt{p})$
- Overall communication complexity: $\Theta(n^2/ \sqrt{p})$